Optimal Metroplex Routing Paradigm for Flexible Flights

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I. Introduction

Air transportation system inefficiencies are found when air traffic demand exceeds system capacity. The resulting flight delays thus increase operating costs of airlines, air traffic controllers’ workload, and the probability of passengers missing connecting flights. Moreover, the air traffic demand is expected to continue its rapid growth in the future. The Federal Aviation Administration estimated in 2007 that the number of passengers is projected to increase by an average of 3% every year until 2025. Most of the traffic demand increase will occur in metropolitan areas where there are two or more large airports. The Joint Planning and Development Office defines this kind of region, with a group of two or more nearby airports for which the arrival and departure operations are highly interdependent, as a metroplex. The New York Metroplex [90 terminal radar approach control facilities (TRACON)], for example, consists of John F. Kennedy International Airport (JFK), LaGuardia airport (LGA), and Newark Liberty International Airport (EWR), as well as several smaller airports. The study in [1] shows that the traffic in most metroplexes has increased significantly over the past years, and the N90 metroplex has the heaviest traffic demand. Future traffic growth will put current facilities under extreme pressure. Therefore, determining how to reduce delays at metroplexes is critical.

The metroplex phenomenon was first studied by Atkins in the San Francisco Bay area and N90 [2]. It was shown that metroplex phenomena would affect the total capacity of the metroplex airport system. In [3], Atkins and Engelland presented a metroplex definition based on several measurable dimensions and proposed an initial framework to study the nature of metroplexes. Wang et al. [4] and Donohue et al. [5] studied the airports in N90 early in 2008, and they considered the status of each airport in terms of the markets served, seat capacity, delays, and other features. McClain and Clarke [6] designed the metric for metroplex clustering analysis. The dependencies and impacts of the three major airports (JFK, LGA, and EWR) in the N90 metroplex were analyzed by Ayyalasomayajula and DeLaurentis in [7,8], where several metrics and models were developed to provide insights for planners to formulate policies and strategies that streamline metroplex operations and mitigate delays.

Although the metroplex concept and its corresponding metrics have been studied by a few researchers, there are very few of them discussing how to increase the metroplex efficiency and reduce the total delay. In this Note, an integrated flexible flight operation simulation platform is introduced, which consists of multiple functional modules such as routing, scheduling, flexible flight selection, and flexible flight candidate plan estimation. Particularly, the routing module is described in detail. The metroplex routing algorithms are implemented by constructing a network data structure from the waypoints (navigational aids, arrival fixes, metering fixes, etc.) and jet routes. Two algorithms are applied with and without the en route traffic congestion constraint, respectively. To the authors’ best knowledge, this work is among the very first integrations of the complete flexible flight operation simulation platform with multiple well-developed NASA- and university-developed tools.

The rest of this Note is organized as follows. The second section, Sec. II, introduces the concepts of the metroplex and flexible flights. The overview of the integration platform and its functional modules is also presented. Section III shows how the network data structure is established and the metroplex routing algorithms with/without traffic congestion constraints are applied. Section IV shows the simulation results of the N90 metroplex, and Sec. V concludes the Note.

II. Overview

A. Metroplex and Flexible Flight

A metroplex is a region with several close airports that share traffic resources such as airspace and ground transportation. More rigorously, a metroplex consists of several close airports with consequential dependencies. Each metroplex includes the airports, the flights and ground traffic between them, the airline companies, and air traffic control services, etc. In this Note, the N90 metroplex is investigated with its three major airports serving the New York City metropolitan area.

A flexible flight is routed toward a metroplex instead of a predetermined destination airport. Unlike the regular flight having a fixed destination airport, when the flexible flight approaches the decision boundary of the destination metroplex, it will receive the instruction of which runway of which airport to land in. The flexible flight operation is based on those passengers who do not have a strong a preference on arriving at a specific destination airport in a metropolitan area. More details about flexible flight operation can be found in [9]. To operate a flexible flight from its takeoff to landing includes two parts. The first part is to route the flexible flight toward its destination metroplex. The second part is when the flexible flight approaches the decision boundary of the metroplex; it will be scheduled to a certain runway in the metroplex according to all the airport conditions. In practice, the decision boundary consists of a series of metering fixes [10].

The concepts of metroplex and flexible flights have two main advantages. First, they allow the air traffic control system to maximize resource utilization (runways in this case) in an otherwise
tightly constrained system. Second, they allow users to experience less delay when accessing a metroplex area.

B. Platform Overview

To simulate the flexible flight operations and evaluate the optimal routing algorithm, an integrated platform is created that consists of several functional modules. The data flow diagram in Fig. 1 illustrates how these modules interact with each other.

Besides the optimal routing module called the linear time-variant optimization module (LTV), there are five major functional modules in Fig. 1. They are ACES (airspace concept evaluation system) [11], the flexible flight selector (FFS), the candidate flight plan estimator, the Multi-Center Traffic Management Advisor (McTMA) scheduler [12], and the airport/runway balancer.

ACES [11] is a NASA software tool that loads in the recorded or optimized flight plans and outputs the high-accuracy simulated aircraft trajectory data with airline identification (ID), computer ID, aircraft type, weight category, origin airport, destination airport, the complete trajectory, estimated time of arrival (ETA), airspeed along the trajectory, etc. The output aircraft trajectory data are classified by FFS into regular flights and flexible flights.

The FFS module was developed by Purdue University, which takes the same set of aircraft trajectory data and designs the flexible flights based on aircraft total travel distance, destination airport, percentage of connecting passengers, and other factors (please see [9] for detail). It categorizes the aircraft trajectory data into regular flights that are fed into McTMA and flexible flights that are loaded into candidate flight plan estimator.

The aircraft trajectory data of flexible flights are sent to the candidate flight plan estimator. Each designated flexible flight only has one recorded trajectory. Based on this existing trajectory, the new candidate flight plans and corresponding ETAs are estimated for the same flexible flight to all the other destination airports inside the metroplex (see [9] for detail).

McTMA [12] takes the regular flights and runs the scheduler to find out the available time slots (holes) for the flexible flights. These holes are then sent to the airport/runway balancer that will decide which destination airport a flexible flight should land in. Then, each flexible flight only has one determined destination airport and one flight plan. These determined flexible flight plans are fed back into the McTMA, where they are scheduled together with the regular flights (Fig. 2).

The airport/runway balancer is a university-developed plug in for the McTMA, which loads in all the candidate flight plans for each flexible flight from the candidate flight plan estimator and only selects one flight plan for every flexible flight. All the selected flexible flight plans are sent back to the McTMA, where they are scheduled.

III. Data Structure and Optimal Routing Algorithms

A. Jetway Network Data Structure

A jetway network data structure is constructed to design the optimal routing algorithm under en route traffic congestion constraint. The data structure is built from navigational aids, fixes, waypoints, and jet routes provided by the Future ATM Concepts Evaluation Tool (FACET), [13] as shown in Fig. 3.

A hash index is used to maintain the data structure, i.e., the nodes (all kinds of waypoints) and edges (jet routes). For ease of demonstration, only the waypoint hash index is described here. The hash key of each waypoint is the rounded positive latitude–longitude pair. In other words, each waypoint’s latitude and longitude are transformed to their absolute values and then rounded into integers. The resulting positive integers are stored in a hash key pair (lat, lon).

From another perspective, applying this kind of hash index is to divide the U.S. airspace into grids. Each grid is a square with four corners for which the latitude and longitude are integers, as shown in Fig. 4a, where i and j are integers. When there is a new waypoint to be added or removed, first calculate its hash key pair and decide which grid it is in. As floor rounding is used to obtain (lat, lon) from the absolute latitude and longitude, if the waypoint exists in the data structure, it should be in the grid with top-left corner coordinates (i, j) = (lat, lon). Therefore, only the existing waypoints (WPs) WP_i, WP_j, and WP_k in this grid (see Fig. 4a) need to be checked instead of checking all the waypoints in the whole data structure. This is an example of the benefit of the hash index, and this method substantially enhances the speed of maintaining waypoint data records.

The hash table H corresponding to the grid in Fig. 4a is kept with a sorted hash key pair, as shown in Fig. 4b. The grids are saved in ascending order of the first integer of the hash key pair, and with the same first integer, the grids are recorded in ascending order of the second integer of the hash key pair. The waypoints in each grid are recorded under the hash key pair (lat, lon), which is also actually the top-left coordinates (i, j) of the corresponding grid. If there is no waypoint in a grid, that grid will not be stored in the hash table.

B. Optimal Routing Algorithms

In this section, the routing algorithm without en route congestion constraint is first described. The edge weight is a great circle distance that is constant. The shortest path algorithm is applied to find the optimal route for each aircraft. Then, the en route congestion constraint is taken into account and the minimum dynamic cost path algorithm is adopted to find the optimal route in a time-variant weighted network. In this research project, the implemented minimum dynamic cost path algorithm is called the LTV optimization module.

1. Optimal Routing Without En Route Congestion Constraint

The network data structure constructed in Sec. III is used for performing the optimal routing algorithm. When the routing algorithm does not take the dynamical traffic congestion into account, the weight of each edge is the great circle distance between two waypoints, which is time invariant. Therefore, the Dijkstra algorithm [14] is adopted to find the optimal route. For regular flights, the algorithm is applied between origin and destination airports. For flexible flights, the routing algorithm terminates at the metroplex decision boundary (metering fix).
2. Linear Time-Variant Optimization Algorithm

To introduce the timely changing en route traffic congestion constraint, the edges intersecting those congested areas will be assigned with additional penalty weights. The time-variant congestion penalty weight will be summed up with the great circle distance weight for each edge and stored in the network data structure. The dynamic congestion information can be retrieved from air traffic controllers. The minimum dynamic cost path algorithm in [15] is implemented to calculate the optimal route with time-variant edge weights.

IV. Simulation

To evaluate the benefits of flexible flight operation and the optimal routing module, actual flight data into the N90 metropolis was tested by the integrated model. The simulation was performed on the historical data of 7 November 2008 at three major airports of N90 (JFK, LGA, and EWR). The Aircraft Situation Display to Industry dataset contains all types of air traffic, such as commercial aircraft, general aviation, and air taxis.

First of all, without the LTV optimization, different percentages of flexible flights were compared in order to study the advantage of flexible flight operation in N90. Zero, 10, 20, 30, 40, and 50% flexible flights were selected by FFS out of the historical flights and directly sent to McTMA scheduling, in which the 0% case with no flexible flights is the evaluation baseline. Figure 5 and Table 1 show how total delay changes according to the percentage of flexible flights. The delay reduction for Newark Liberty International Airport (KEWR), John F. Kennedy International Airport (KJFK) and La Guardia Airport (KLGA) is listed. As the percentage of flexible flights (FFs) increases, the total delay of the three major airports decreases almost linearly. With the flexible flight operation concept, the airport/runway balancer reduces the total delay by rescheduling a flexible flight to the less-congested airport. Compared to the baseline case (no flexible flight), the total delay reduction reaches 29.4% with 50% flexible flights.

![Fig. 3 Waypoints and jet routes provided by FACET.](image)

![Fig. 4 Representations of a) hash-indexed grids enhancing the maintenance of waypoints, and b) waypoints stored in a sorted hash table H.](image)

![Fig. 5 Total delays for various flexible flight percentages into N90 on 7 November 2008.](image)

<table>
<thead>
<tr>
<th>KEWR</th>
<th>KJFK</th>
<th>KLGA</th>
<th>Total</th>
<th>Reduction, %</th>
</tr>
</thead>
<tbody>
<tr>
<td>Nominal (no FF)</td>
<td>610.1</td>
<td>948.6</td>
<td>466.3</td>
<td>2025.0</td>
</tr>
<tr>
<td>10%</td>
<td>586.6</td>
<td>1001.4</td>
<td>316.2</td>
<td>1904.2</td>
</tr>
<tr>
<td>20%</td>
<td>547.9</td>
<td>975.3</td>
<td>313.3</td>
<td>1836.5</td>
</tr>
<tr>
<td>30%</td>
<td>414.7</td>
<td>967.0</td>
<td>322.4</td>
<td>1704.2</td>
</tr>
<tr>
<td>40%</td>
<td>377.4</td>
<td>825.5</td>
<td>406.2</td>
<td>1609.1</td>
</tr>
<tr>
<td>50%</td>
<td>389.3</td>
<td>586.4</td>
<td>453.4</td>
<td>1429.1</td>
</tr>
</tbody>
</table>

- **Fig. 4** Representations of a) hash-indexed grids enhancing the maintenance of waypoints, and b) waypoints stored in a sorted hash table H.
- **Fig. 5** Total delays for various flexible flight percentages into N90 on 7 November 2008.
The result of the second study is shown in Fig. 6, in which, with LTV optimal routing and flexible flight operation (FFO) on/off, the comparison of four different combinations is illustrated. Without flexible flight operation, LTV reduced the total delay at the three airports from 2025 to 1830 min (9.5% reduction) as compared to the baseline case (no LTV optimization and no flexible flight operation). Without the LTV, applying flexible flight operation to 30% of the flights in historical data (30% are flexible flights, and 70% are regular flights) gives about 300 min total delay reduction because the flexible flight operation (FFO) well balanced the runway resource. In summary, combining LTV optimization and flexible flight operations has the largest delay reduction (536 min reduction from 2025 min). The simulation in Fig. 6 shows that both the LTV optimization module and flexible flight operation are beneficial to the metroplex in terms of reducing total delay.

V. Conclusions

In this Note, the concepts of flexible flight and metroplex are presented. To increase the efficiency and reduce the total delay of the metroplex airports, an integrated flexible flight simulation platform is developed, which includes multiple functional modules. The major contributor for total delay reduction is the optimal routing module based on the jetway network data structure. All the other modules are briefly introduced. The numerical result shows the benefits of both the flexible flight operation and the optimal routing module. The developed integrated simulation system can be considered as the framework of further metroplex studies.

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References