# New first-order methods in modern/classical settings

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Part 1

# Sharp exact penalty formulations in signal recovery

Joint work with Lijun Ding

Ding, Wang

• Motivation: Sparse recovery and low-rank covariance estimation

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  - $\bullet \longrightarrow \mathsf{Abstract signal recovery problem}$

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  - $\longrightarrow$  Abstract signal recovery problem
- A new formulation of the abstract problem that is sharp
  - Better robustness guarantees, faster algorithms
- Numerical results

# Motivation: Sparse recovery and covariance estimation

• **Recovery task**: Recover  $x^{\sharp} \in \mathbb{R}^n$  from  $A \in \mathbb{R}^{m \times n}$ ,  $b = Ax^{\sharp}$ 

Related: Candes and Tao [2005], Recht et al. [2010], Candès et al. [2013]

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- Convex optimization approach: In this regime,  $x^{\sharp}$  is unique minimizer of

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  - Another convex problem?

# A sharp penalty formulation

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Related: Beck and Teboulle [2009], Tibshirani [1996]

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$$\min_{x \in V} \left\{ f(x) : \begin{array}{l} \mathcal{A}(x) = b \\ x \in K \end{array} \right\}$$

• Penalty formulation: let  $r \asymp \sqrt{k}$  be a penalty parameter

 $F(x) \coloneqq f(x) + \frac{r \left\| \mathcal{A}(x) - b \right\|_1 + 2 \operatorname{dist}_1(x, K)}{r \left\| \mathcal{A}(x) - b \right\|_1 + 2 \operatorname{dist}_1(x, K)}$ 

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$$\|Ax - b\|_2^2$$
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$$F(x) - F(x^{\sharp}) \ge \mu \left\| x - x^{\sharp} \right\|_{1}, \qquad \forall x \in V$$

and L-Lipschitz in the  $\ell_1$  norm with  $L \asymp \sqrt{k}$ 

$$|F(x) - F(y)| \le L ||x - y||_1, \quad \forall x, y.$$

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#### Corollary (Robustness)

Let  $\tilde{x}$  be an  $\epsilon$  minimizer of  $\tilde{F}$ .

• (to small noise)  $\tilde{x}$  satisfies  $\|\tilde{x} - x^{\sharp}\|_{1} \lesssim \frac{\sqrt{k}}{\mu} \|\delta\|_{1} + \frac{\epsilon}{\mu}$ 

• (to sparse noise) If 
$$rac{|\mathrm{supp}(\delta)|}{m} \lesssim 1/\sqrt{k}$$
, then  $\left\|\tilde{x} - x^{\sharp}\right\|_1 \lesssim rac{\epsilon}{\mu}$ 

# Algorithms for minimizing *F*



## Corollary (Algorithms)

Restarted mirror descent (RMD) algorithm produces an  $\epsilon$ -optimal solution to F in

$$O\left(\frac{k}{\mu^2}\log(n)\log(\epsilon^{-1})\right)$$

iterations of the mirror descent update.

Related: Polyak [1969], Roulet and d'Aspremont [2017], Yang and Lin [2018], Renegar and Grimmer [2022]

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- If  $\mu$  is not known, extra  $\log(\epsilon^{-1})$  factor

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• Suppose we run MD from  $x_0$  for t iterations with step size  $\eta$  and mirror map

$$h(x) \approx \frac{1}{2} \|x - x_0\|_1^2$$

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• Applying sharpness  $\longrightarrow$ 

$$F(y) - F(x^{\sharp}) \le \frac{1}{2} \left( F(x_0) - F(x^{\sharp}) \right)$$

after  $\asymp \frac{L^2}{\mu^2} \ln n$  iterations

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# Numerical experiments

#### **Restarted mirror descent**

• Let *T* be statistical threshold for sparse recovery, low-rank matrix sensing (covariance estimation without PSD constraint), and phase retrieval (covariance estimation with k = 1)



#### Restarted mirror descent vs. Polyak subgradient

• Polyak subgradient converges linearly on sharp Lipschitz functions in  $\ell_2$  norm



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# **Questions?**

# Part 2 $O(1/T^{1.02449})$ Convergence of long-step gradient descent

Joint work with Benjamin Grimmer, Kevin Shu

Grimmer, Shu, Wang

Accelerated convergence rates for gradient descent

• Preview of results (better guarantees for smooth convex minimization)

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# **Preview of results**

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$$x_1 = x_0 - h_0 \nabla f(x_0)$$
  $x_2 = x_1 - h_1 \nabla f(x_1)$  ...  
 $x_{i+1} = x_i - h_i \nabla f(x_i)$ 

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• Goal: pick steplength sequence  $(h_0, h_1, \dots)$  to maximize convergence rate

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Today: a per-iteration analysis is too short-sighted

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Today: a per-iteration analysis is too short-sighted

• Optimal rates for first-order methods: Accelerated gradient descent

$$f(x_T) - f(x^\star) \le \frac{2}{T^2}$$

Note: this is not a gradient descent-style algorithm

• Consider 
$$h = 0.99 \times \left( \left| \frac{3}{2}, 5, \frac{3}{2} \right|, \left| \frac{3}{2}, 5, \frac{3}{2} \right|, \ldots \right)$$

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#### Taking larger steps: breaking some intuitions

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$$h^{(0)} = (1), \quad h^{(1)} = \left(\frac{3}{2}, 5, \frac{3}{2}\right), \quad h^{(2)} = \left(\frac{3}{2}, 1 + \sqrt{2}, \sqrt{2}, 7 + 4\sqrt{2}, \sqrt{2}, 1 + \sqrt{2}, \frac{3}{2}\right)$$

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- We construct steplength blocks  $h^{(k)} \in \mathbb{R}^{2^{k+1}-1}$  that can be scaled down to guarantee descent

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- Longer patterns have increasingly fast convergence rates
- $\operatorname{avg}(h^{(k)})$  is exponential in k

# A closer look at $h^{(k)}$



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- $\alpha_i$  picked so that  $\prod_{\text{stepsizes}}(\text{stepsize} 1) = 1$

# Numerical comparison of $h^{(k)}$



#### Accelerated convergence for gradient descent-style algorithms

## Theorem

#### Suppose

$$h = \frac{1}{2} \left( \boxed{h^{(0)}, \dots, h^{(0)}}, \boxed{h^{(1)}, \dots, h^{(1)}}, \dots, \boxed{h^{(k)}, \dots, h^{(k)}}, \dots \right)$$

where each  $h^{(k)}$  is repeated  $\approx c^k$  times. Then

$$\left(\min_{t \le T} f(x_t)\right) - f(x^{\star}) = O\left(\frac{1}{T^{1.02449}}\right)$$

# Why should we expect this?

AKA some recent work in the area

• **Question**: Suppose we have a candidate  $h = (h_0, h_1, ..., h_{T-1})$ . What is the worst case function? (Smoothness 1, initial distance 1, initial suboptimality  $\delta$ )

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$$p_{h}(\delta) \coloneqq \max_{x_{0}, x^{\star}, f} \begin{cases} f \text{ is convex, 1-smooth} \\ \|x_{0} - x^{\star}\|^{2} \leq 1 \\ f(x_{T}) - f(x^{\star}) : f(x_{0}) - f(x^{\star}) \leq \delta \\ \nabla f(x^{\star}) = 0 \\ x_{i+1} = x_{i} - h_{i} \nabla f(x_{i}) \end{cases}$$

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  - The SDP relaxation of this nonconvex quadratic program is exact!

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#### • Take-aways:

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- Now, how to design h?

$$\min_{h=(h_0,\ldots,h_{T-1})} p_h(\delta)$$

• Das Gupta et al. [2023]: Complex branch-and-bound scheme for  $T \in [1, \dots, 50]$ 

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3.5 15 30  $h^{lopt}$ 10 foot 20 foot tdo 2.5 . 2 10 2 1.5 . 10 20 25 0 10 20 40 50 0 20 Iteration number Iteration number Iteration number Iteration number (a) N = 5(b) N = 10(c) N = 25(d) N = 50

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 $f(x_{50}) \leq 0.002 \approx$  factor of 5 faster than  $\frac{1}{2T}$ 

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- Our work: analytically construct solution for all  $\delta$ , with  $p_h(\delta)$  small

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- Loshchilov and Hutter [2016], Smith [2015], Smith and Topin [2017]: Nonconvex, smooth minimization in neural networks

# **Conceptual contributions**

• Intuition: Long flat regions of small slope is the worst case

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- Suppose  $\delta > 0$  small and consider

$$f(x) = \begin{cases} -\delta x - \delta^2/2 & \text{if } x \leq -\delta \\ \frac{1}{2}x^2 & \text{if } x \geq -\delta \end{cases}$$

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• Then,  $f(x_0) = \delta$  and  $f(x_T) = \delta - \delta^2 \sum h_i$ 

• Thus, for  $\delta > 0$  small,  $p_h(\delta) \ge \delta - \delta^2 \sum h_i$ 

• We say a steplength *block*  $h = (h_0, h_1, \dots, h_{T-1})$  is  $\Delta$ -straightforward if

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- We show that  $\operatorname{avg}(h^{(k)})$  is exponentially large in k,  $\Delta^{(k)}$  is  $\geq$  exponentially small  $\longrightarrow$  accelerated convergence rates

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• This becomes a nonlinear SDP but can be "reformulated" into a regular SDP if we consider "limiting behavior as  $\Delta \rightarrow 0$ ", at which point we can attempt to certify  $\Delta$ -straightforwardness computationally

# Computer assisted design/proofs

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#### **Pictures of process**



Grimmer, Shu, Wang

Accelerated convergence rates for gradient descent

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# **Questions?**

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